

**Solution Concepts 3**  
**Nash equilibrium in pure strategies**  
Watson §9-§10, pages 89-100

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## Rationalizability vs equilibrium

- By assuming that there is common knowledge of rationality, we have concluded that players will choose rationalizable strategies
- This prediction has two criticisms:
  - ① In most cases it is not very informative
  - ② It allows players to have erroneous beliefs
- By assuming that players beliefs are correct (i.e. if player 1 has beliefs  $\theta_2$  about 2's behavior then 2 makes choices according to  $\sigma_2 = \theta_2$ ) we obtain different notions of equilibrium
- In this slides we only consider *Nash equilibrium in pure strategies*

# Correct beliefs

- *Why would we assume that players have correct beliefs?*
  - ① Communication.– If players communicate with each other prior to playing the game they might agree to follow some strategies
  - ② Learning.– If players interact repeatedly they might learn from experience how to predict their opponents behavior
  - ③ Adaptation.– If players follow simple adaptive rules, behavior can also converge to something that looks like an equilibrium
  - ④ Institutions.– Institutions/mediators might help to coordinate players expectations
  - ⑤ Focal points.– Some rationalizable strategies might be justifiable by simple logical arguments

# Nash equilibrium in pure strategies

## Communication and self-enforcing agreements

- Suppose that the players gather to discuss and agree on playing according to some strategy profile  $s \in S$  specifying a pure strategy for each player (no mixing for now)
- After that, players go different ways and they choose strategies simultaneously and independently
- Suppose that player  $i$  thinks that his/her opponents will not deviate from the agreed strategy profile, i.e. that they will choose the strategies in  $s_{-i}$
- Then  $i$  will be willing to choose strategy  $s_i$  if and only if it is a best response to  $s_{-i}$ , i.e. if and only if  $s_i \in BR_i(s_{-i})$
- In this case  $i$  can not **strictly** benefit from **unilaterally** deviating from the intended strategy profile
- If no players have strict incentives to deviate **unilaterally** then we say that  $s$  is a Nash equilibrium in pure strategies

# Nash equilibrium in pure strategies

## Definition

Given a strategic form game, a Nash equilibrium *in pure strategies* is a strategy profile  $s \in S$  such that no player can **strictly** gain from deviating **unilaterally**, i.e. such that:

$$u_i(s_i, s_{-i}) \geq u_i(s'_i, s_{-i})$$

for every player  $i$  and every alternative strategy  $s'_i \in S_i$

- Equivalently, a Nash equilibrium is a profile of strategies which are best responses to each other, i.e. a strategy profile  $s \in S$  such that  $s_i \in BR_i(s_{-i})$  for every player  $i$
- In a two player game (represented by a payoff matrix) a pair of strategies is a Nash equilibrium if player 1 is maximizing his/her payoff along the corresponding *column* and player 2 is maximizing his/her payoff along the corresponding *row*

## Example: Battle of the Sexes

Nash equilibria

	Football	Opera
Football	<u>5</u> , <u>1</u>	0, 0
Opera	0, 0	<u>1</u> , <u>5</u>

- To find Nash equilibrium of a finite game one can start by highlighting the best response payoffs for each player
- If a cell in the matrix has all payoffs highlighted then it is a Nash equilibrium

# Rationalizability vs Nash equilibrium

- If we assume that:
  - ① Players are rational
  - ② Players are making deterministic choices (no mixed strategies)
  - ③ Players have correct beliefs about their opponents' behavior (they know what their opponents are going to choose)
- then we can predict that they will play some Nash equilibrium
- Nash equilibria are *joint* predictions specifying strategies for *all* players
- Rationalizability makes *individual* predictions for *each* player

## Theorem

*Every strategy in a Nash equilibrium is rationalizable*

## Theorem

*If there is a unique rationalizable strategy for each player, then these strategies conform a Nash equilibrium*

# Example: A $4 \times 4$ game

Best responses

	a	b	c	d
w	0, <u>7</u>	2, 5	<u>7</u> , 0	0, 1
x	5, 2	<u>3</u> , <u>3</u>	5, 2	0, 1
y	<u>7</u> , 0	2, 5	0, <u>7</u>	0, 1
z	0, <u>0</u>	0, -2	0, <u>0</u>	<u>10</u> , -1

# Example: A $4 \times 4$ game

Nash equilibrium and rationalizable strategies

	a	b	c	d
w	0, <u>7</u>	2, 5	<u>7</u> , 0	0, 1
x	5, 2	<u>3</u> , <u>3</u>	5, 2	0, 1
y	<u>7</u> , 0	2, 5	0, <u>7</u>	0, 1
z	<del>0, <u>0</u></del>	<del>0, -2</del>	<del>0, <u>0</u></del>	<del><u>10</u>, -1</del>

# Example: classic $2 \times 2$ examples

Best responses

	Full	Empty
Full	3, 3	0, <u>5</u>
Empty	<u>5</u> , 0	<u>2</u> , <u>2</u>

	Continue	Swerve
Continue	0, 0	<u>5</u> , <u>1</u>
Swerve	<u>1</u> , <u>5</u>	2, 2

	GCS	PS
GCS	<u>1</u> , <u>1</u>	0, 0
PS	0, 0	<u>2</u> , <u>2</u>

	Press	Don't press
Press	3, 1	<u>0</u> , <u>5</u>
Don't press	<u>6</u> , -2	-1, <u>-1</u>

# Example: classic $2 \times 2$ examples

Nash and rationalizability

	Full	Empty
Full	<del>3, 3</del>	<del>0, <u>5</u></del>
Empty	<u>5</u> , 0	<u>2</u> , <u>2</u>

	Continue	Swerve
Continue	0, 0	<u>5</u> , <u>1</u>
Swerve	<u>1</u> , <u>5</u>	2, 2

	GCS	PS
GCS	<u>1</u> , <u>1</u>	0, 0
PS	0, 0	<u>2</u> , <u>2</u>

	Press	Don't press
Press	3, 1	<u>0</u> , <u>5</u>
Don't press	<u>6</u> , -2	<del>1, <u>1</u></del>

## Example: rock paper scissors

Not every game has a Nash equilibrium in pure strategies

	Rock	Paper	Scissors
Rock	0, 0	-1, <u>1</u>	<u>1</u> , -1
Paper	<u>1</u> , -1	0, 0	-1, <u>1</u>
Scissors	-1, <u>1</u>	<u>1</u> , -1	0, 0

# Example: Cournot competition

## Best responses

- Consider a Cournot duopoly game with two firms 1 and 2 choosing quantities  $q_1, q_2 \in [0, 50]$ , with constant marginal costs  $c = 10$  and inverse demand function:

$$P(q_1, q_2) = 100 - q_1 - q_2$$

- Payoffs are given by:

$$u_1(q_1, q_2) = (90 - q_2 - q_1)q_1 \quad u_2(q_1, q_2) = (90 - q_1 - q_2)q_2$$

- Best responses to pure strategies are given by:

$$BR_1(q_2) = 45 - \frac{1}{2}q_2 \quad BR_2(q_1) = 45 - \frac{1}{2}q_1$$

# Example: Cournot competition

## Nash equilibria

- A pure strategy Nash equilibrium for this Cournot example is a pair of quantities  $(q_1, q_2)$  that are mutual best responses, i.e. such that:

$$q_1 = BR_1(q_2) \quad q_2 = BR_2(q_1)$$

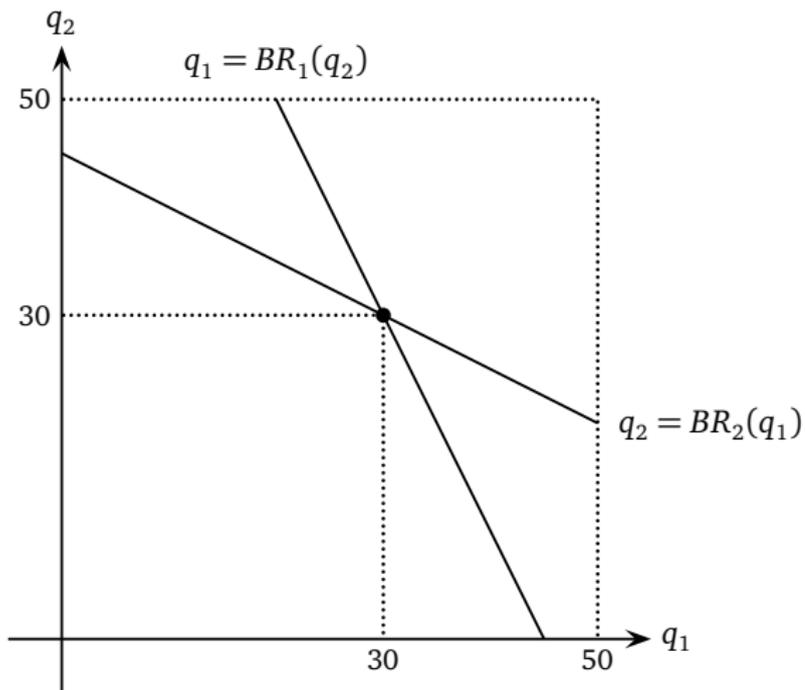
- Using our formula for best responses this is equivalent to:

$$\begin{aligned}q_1 &= 45 - \frac{1}{2}q_2 & q_2 &= 45 - \frac{1}{2}q_1 \\ \Rightarrow q_2 &= 45 - \frac{1}{2} \left( 45 - \frac{1}{2}q_2 \right) = 45 - 22.5 + \frac{1}{4}q_2 = 22.5 + \frac{1}{4}q_2 \\ \Rightarrow \frac{5}{4}q_2 &= 22.5 & \Rightarrow q_2 &= \frac{4 \cdot 22.5}{5} = 30 \\ \Rightarrow q_1 &= 45 - \frac{1}{2}30 = 45 - 15 = 30\end{aligned}$$

- So the game has a unique Nash equilibrium in pure strategies:  $(30, 30)$
- Recall that this was the *unique* rationalizable strategy profile

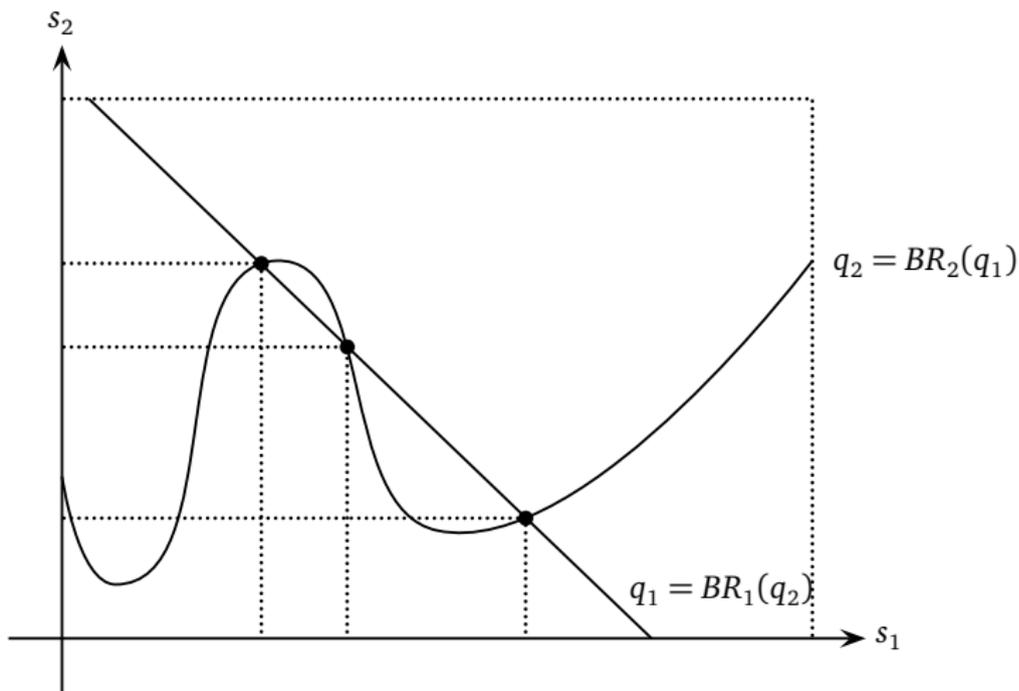
# Example: Cournot competition

Nash equilibrium



# Example: a continuous two player game

Best responses and Nash equilibrium



# Example: location game

Nash equilibrium

	1	2	3	4	5	6	7
1	35, 35	10, <u>60</u>	15, 55	20, 50	25, 45	30, 40	35, 35
2	<u>60</u> , 10	35, 35	20, <u>50</u>	25, 45	30, 40	35, 35	40, 30
3	55, 15	<u>50</u> , 20	35, 35	30, <u>40</u>	35, 35	40, 30	45, 25
4	50, 20	45, 25	<u>40</u> , 30	<u>35</u> , <u>35</u>	<u>40</u> , 30	45, 25	50, 20
5	45, 25	40, 30	35, 35	30, <u>40</u>	35, 35	<u>50</u> , 20	55, 15
6	40, 30	35, 35	30, 40	25, 45	20, <u>50</u>	35, 35	<u>60</u> , 10
7	35, 35	30, 40	25, 45	20, 50	15, 55	10, <u>60</u>	35, 35