

# Solution Concepts 4

## Nash equilibrium in mixed strategies

Watson §11, pages 123-128

*Bruno Salcedo*

The Pennsylvania State University

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# Mixing strategies

- In a strictly competitive situation players have incentives to prevent their opponents from predicting their choices
  - Examples: rock paper scissors, military tactics, poker
- One way of remaining “unpredictable” is to randomize your choices

## Definition

A mixed strategy for player  $i$  is a probability distribution  $\sigma_i$  over his/her strategies

- See the slides on dominance and best responses (S4) or section §5 in the textbook for more details.
- We don't think of actual explicit randomization (eg rolling a dice to make a choice) but rather implicit randomization (eg basing your choices on “feelings” or unpredictable introspective processes)
- We use the adjective “pure” to talk about non-mixed strategies. A pure strategy is equivalent to the mixed strategy that plays it for sure

# Nash equilibrium in mixed strategies

- When players randomize, we can compute expected utility:

$$\begin{aligned} U_i(\sigma_i, \sigma_{-i}) &= \mathbb{E} \left[ u_i(s_i, s_{-i}) \mid \sigma_i, \sigma_{-i} \right] \\ &= U_i(\sigma_i, \sigma_{-i}) = \sum_{s_i \in S_i} \sum_{s_{-i} \in S_{-i}} u_i(s_i, s_{-i}) \quad (\text{for finite games}) \end{aligned}$$

- The notions of rationality, rationalizability, best responses and Nash equilibrium remain unchanged

## Definition

Given a strategic form game, a Nash equilibrium is a (pure or mixed) strategy profile  $\sigma$  such that no player can **strictly** gain from deviating **unilaterally**, i.e. such that:

$$U_i(\sigma_i, \sigma_{-i}) \geq U_i(\sigma'_i, \sigma_{-i})$$

for every player  $i$  and every alternative strategy  $\sigma'_i$

## Example: Rock Paper Scissors

	Rock	Paper	Scissors
Rock	0, 0	-1, 1	1, -1
Paper	1, -1	0, 0	-1, 1
Scissors	-1, 1	1, -1	0, 0

- Claim: both players randomizing according to  $(1/3, 1/3, 1/3)$  is a Nash equilibrium
- If a player uses this strategy his/her opponent's expected payoff *for any strategy* is 0
- Thus there are no incentives to deviate unilaterally

# Computing equilibria in mixed strategies

## Theorem

*If a mixed strategy  $\sigma_i$  is a best response to  $\sigma_{-i}$  then so are all the strategies that are mixed with positive probability*

- This means that, if a player is willing to randomize, it must be the case that he/she is *indifferent* between all the strategies over which he is randomizing
- To find Nash equilibria in mixed strategies we do the following:
  - ① “Guess” the pure strategies that will be mixed (start by eliminating strategies that are not rationalizable)
  - ② For each player  $i$ , look for a mixed strategy for  $-i$  that makes  $i$  be indifferent between the strategies that he/she is mixing

## Example: A $2 \times 2$ game

Row's expected utility

		Col	
		L [ $p$ ]	R [ $1-p$ ]
Row	U [ $q$ ]	3, 3	5, 8
	D [ $1-q$ ]	1, 2	6, 1

- Given  $p$ , row's expected utility for each pure strategy is:

$$U_1(U, p) = 3p + 5(1-p) = 5 - 2p$$

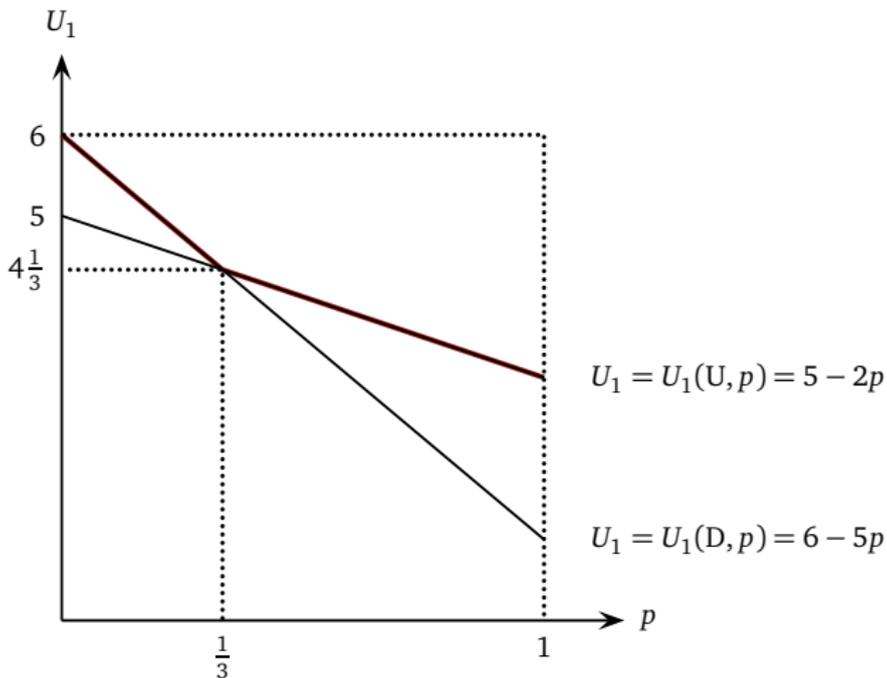
$$U_1(D, p) = 1p + 6(1-p) = 6 - 5p$$

- Row is thus indifferent between U and D if and only if:

$$U_1(U, p) = U_1(D, p) \iff 5 - 2p = 6 - 5p \iff p = \frac{1}{3}$$

# Example: A $2 \times 2$ game

Row's best responses



## Example: A $2 \times 2$ game

		Col	
		L [ $p$ ]	R [ $1-p$ ]
Row	U [ $q$ ]	3, 3	5, 8
	D [ $1-q$ ]	1, 2	6, 1

- Given  $q$ , Col's expected utility for each pure strategy is:

$$U_2(L, q) = 3q + 2(1 - q) = 2 - q$$

$$U_2(R, q) = 8q + 1(1 - q) = 7q - 1$$

- Col is thus indifferent between L and R if and only if:

$$U_2(L, q) = U_2(R, q) \iff 2 - q = 7q - 1 \iff q = \frac{1}{6}$$

## Example: A $2 \times 2$ game

		Col	
		L [ $p$ ]	R [ $1-p$ ]
Row	U [ $q$ ]	3, 3	5, 8
	D [ $1-q$ ]	1, 2	6, 1

- We then have found a mixed equilibrium in pure strategies:

$$\sigma_1 = \left( \frac{1}{6}, \frac{5}{6} \right)$$

$$\sigma_2 = \left( \frac{1}{3}, \frac{2}{3} \right)$$

## Why bother making opponent be indifferent?

- It might not seem intuitive that a player randomizes with the exact probabilities that make his/her opponent be indifferent.
- Recall: *making an opponent indifferent is not the intention of the player*, the player simply wants to maximize his expected utility
- The definition and motivation of Nash equilibrium is only that players want to maximize their expected utility, and their beliefs are in equilibrium (there are no profitable unilateral deviations)
- The fact that the corresponding strategies must make players indifferent is a result

## Example: A $4 \times 4$ game

	a	b	c	d
w	0, 9	0, 4	0, 2	0, 6
x	2, 1	9, 3	1, 7	2, 2
y	7, 1	0, 0	3, 5	0, 2
z	2, 1	1, 8	4, 0	1, 4

- Using iterated dominance we end up with a  $2 \times 2$  game
- Let  $p$  be the probability of  $b$  and  $1 - p$  the probability of  $c$ , for indifference we must have:

$$9p + (1 - p) = p + 4(1 - p) \quad \Leftrightarrow \quad p = \frac{3}{11}$$

- Let  $q$  be the probability of  $x$  and  $1 - q$  the probability of  $z$ , for indifference we must have:

$$3q + 8(1 - q) = 7q + 0(1 - q) \quad \Leftrightarrow \quad q = \frac{2}{3}$$

# Existence of equilibrium

## Theorem

Every **finite** strategic form game has **at least** one Nash equilibrium

## Theorem

Generically, finite strategic form games have an odd number of Nash equilibria

# Example: A $2 \times 2$ game

Existence of equilibria

